

## CounterExample Guided Inductive Synthesis modulo Theories

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Program synthesis is hard.

- Extension to CEGIS framework
- Uses general counterexamples and candidates
- Avoids enumerating search space
- Can synthesize programs that elude other solvers



















### Safety invariant

 $init(x) \iff x = 0$  $trans(x, x') \iff x' = x + 1$ 

#### find *inv*(*x*) such that:

 $init(x) \implies inv(x)$ 

 $inv(x) \land (x < 1000) \land trans(x, x') \implies inv(x')$ 

 $inv(x) \land \neg(x < 1000) \implies (x < 1005) \land (x > 5)$ 

### Safety invariant

 $init(x) \iff x = 0$  $trans(x, x') \iff x' = x + 1$ 

$$inv(x) = (4 < x) \land (x < 1003)$$

Target:  
$$inv(x) = (4 < x) \land (x < 1003)$$



And so on ...

# Can we ask more general questions?



# Can we give more general answers?

## More general questions More general answers













### Generalize





### Deduction



is there a value for v that makes (x < v) a valid invariant



### First order solver

Solves 1st order formula with:

- Arbitrary propositional structure
- 1 quantifier alternation

### CEGIS(T) - SMT



### CEGIS(T) - SMT

$$\exists v \forall x . \sigma(P^*[v], x) \land (v < c) \qquad \exists v \forall x . \sigma(P^*[v], x) \land (v > c)$$

$$\neg P^*[v] \qquad v > c \qquad v < c \qquad P^*[v]$$
**BLOCK** CONSTRAINT CONSTRAINT SOLUTION

Target:  
$$inv(x) = (4 < x) \land (x < 1003)$$

$$\exists v \forall x \, . \, \sigma(P^*[v], x) \land (v < c) \qquad \exists v \forall x \, . \, \sigma(P^*[v], x) \land (v > c)$$



Target:  
$$inv(x) = (4 < x) \land (x < 1003)$$

$$\exists v \forall x . \sigma(P^*[v], x) \land (v < 95) \left( \exists v \forall x . \sigma(P^*[v], x) \land (v > 95) \right)$$







Target:  

$$inv(x) = (4 < x) \land (x < 1003)$$
UNSAT
$$\exists v \forall x . \sigma(P^*[v], x) \land (v < 95) \qquad \exists v \forall x . \sigma(P^*[v], x) \land (v > 95)$$

$$\neg P^*[v] \qquad v > 95 \qquad v < 95 \qquad P^*[v]$$
BLOCK
$$P^*[v] \qquad v > 95 \qquad v < 95 \qquad P^*[v]$$
Solution

Target:  
$$inv(x) = (4 < x) \land (x < 1003)$$

$$\exists v \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \qquad \exists v \forall x . \sigma(P^*[v], x) \land (v_1 > 95) \qquad \exists v \forall x . \sigma(P^*[v], x) \land (v_1 > 95) \qquad \forall x . \sigma(P^*[v], x) \land (v_1 > 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 > 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 > 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 > 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 > 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 > 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1 < 95) \ \forall x . \sigma(P^*[v], x) \land (v_1$$

$$P^{*} = (10 < x) \land (x < 95)$$

$$P^{*}[v] = (v_{0} < x) \land (x < v_{1})$$

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Target:inv(x) = (4 < x) \land (x < 1003)TIMEOUTTIMEOUT
$$\exists v \forall x . \sigma(P^*[v], x) \land (v_1 < 95)$$
 $\exists v \forall x . \sigma(P^*[v], x) \land (v_1 > 95)$ 



$$\neg P^*[v] \qquad v > 95 \qquad v < 95 \qquad P^*[v]$$
BLOCK CONSTRAINT CONSTRAINT SOLUTION

### Experiments

Benchmarks:

- Bitvectors
- Syntax-guided Synthesis competition (without the syntax)
- Loop invariants
- Danger invariants

Solvers:

- CVC4
- EUSolver, E3Solver, LoopInvGen bitvectors with no grammar unsupported

### **Experiments**





CEGIS(T) solves program synthesis via 1<sup>st</sup> order solvers that support quantifiers:

Enables use of existing solvers

Algorithmic insights:

- verify generalized candidate solutions
- return generalized counterexamples

#### www.cprover.org/synthesis



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